

# How Minds Work Sparse Distributed Memory

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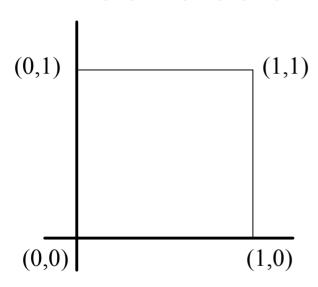
### **Boolean Geometry**

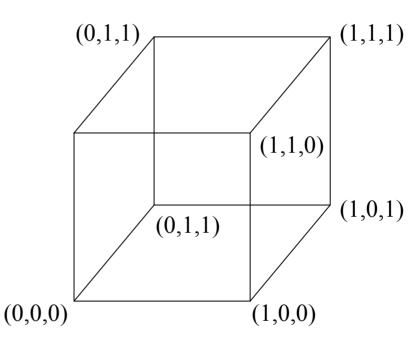
- The geometry of Boolean space
- Boolean space of dimension n— the set of all Boolean vectors of length n
- One dimensional Boolean space— {(0), (1)}
- Two dimensional Boolean space— {(0,0), (0,1), (1,0), (1,1)}



### **Boolean Spaces**

Two dimensional





Three dimensional



## Cardinality of Boolean Spaces

- B<sub>1</sub> has 2<sup>1</sup>=2 elements
- B<sub>2</sub> has 2<sup>2</sup>=4 elements
- B<sub>3</sub> has 2<sup>3</sup>=8 elements
- B<sub>4</sub> has 2<sup>4</sup>=16 elements
- •
- B<sub>n</sub> has 2<sup>n</sup> elements

B<sub>1000</sub> has 2<sup>1000</sup> elements

More than the number of atoms in the universe!



## Hamming Distance

- Number of dimensions in which two Boolean vectors differ
- d((0,1,0,1,0,1,1), (1,1,0,0,0,1,0)) = 3
- 99.9999% of B<sub>1000</sub> lies between 422
   & 578 from a given vector.
- Almost all of a Boolean space is far from any given vector
- Every Boolean space is thinly populated



### **Boolean Spheres**

- Locus of points at some fixed distance, the radius, from its center
- $O(r,x) = \{ y \mid d(x,y) \le r \}$
- For r ≤ n/2 most of the points in O(r,x) lie close to its boundary.



### SDM as Memory

- Random (vs sequential access)
  - Retrieve in equal time from any location
- Content addressable
  - Find complete contents from a part
- Associative
  - Find contents similar to a cue



### Addresses in SDM

- Addresses Boolean vectors of length 1000
- Address space =  $B_{1000}$
- Too enormous to ever implement
- Each dimension a feature, either on (1) or off (0)
- 1000 not many features



### Hard Locations in SDM

- Choose 2<sup>20</sup> (~1,000,000) hard locations
- Uniformly distributed in address space
- 2<sup>20</sup> hard locations out of 2<sup>1000</sup> locations, ratio is 1/2<sup>980</sup> very sparse indeed
- median distance from random location to nearest hard location is 424
- Hard locations are certainly sparse



#### Counters

- Each hard location has 1000 counters
- Each counter has range -40 to 40
- Takes about a gigabyte of memory
- Writing a 1 to a counter increments it;
   writing a 0 decrements it
- Write to a hard location
   — write each
   coordinate to the corresponding counter



### Access Sphere

- Access sphere at some location x sphere of radius 451 centered at x
- Contains about 1000 hard locations
- To write to a location x write to each hard location in its access sphere
- Distributed representation
- Hence Sparse Distributed Memory



### Reading from a Hard Location

- If the ith counter of the hard location is
  - Positive, put a 1 in the ith dimension
  - Negitive, put a 0 in the ith dimension
- This is majority rule at each dimension
- A Boolean vector of the right dimension results
- It may differ from any previously written



### Reading from any Location

- Find the access circle centered at the given location
- Read at each hard location in the circle
- Majority rule over these reading
- Iterate using the result as a new location
- Stop if the itteration stabilizes



### Retrieval

- Items read in (with themselves as address) can be reconstructed
- Iterated reading allows reconstruction from a partial or noisy cue
- Reconstructions may not be exact
- Interference affect occur



### Psychological Effects

- Knowing that one does or doesn't know
- Tip of the tongue feeling



### Readings

- Kanerva, P. 1988. Sparse Distributed
   Memory. Cambridge MA: The MIT Press.
- Franklin, S. 1995. Artificial Minds.
   Cambridge MA: MIT Press



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